Deduction and Computation through Deduction Modulo

Olivier Hermant

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- It has been forgotten by the formalization of the mathematics.
- reborn with informatics: rewriting rules.
- we need a balance between deduction steps and computation steps.

Deduction systems: the logical framework

- first-order logic: function and predicate symbols, logical connectors: $\wedge, \vee, \Rightarrow, \neg$, and quantifiers $\forall, \exists$.

$$\begin{gathered} Even(0) \\ \forall n(Even(n) \Rightarrow Odd(n+1)) \\ \forall n(Odd(n) \Rightarrow Even(n+1)) \end{gathered}$$

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$$\begin{array}{c} \mathit{Even}(0) \\ \forall n(\mathit{Even}(n) \Rightarrow \mathit{Odd}(n+1)) \\ \forall n(\mathit{Odd}(n) \Rightarrow \mathit{Even}(n+1)) \end{array}$$

- a sequent :

$$\overbrace{\Gamma}^{\text{hyp.}} \vdash \overbrace{A}^{\text{conc.}}$$

- rules to form them: sequent calculus (or natural deduction)
- framework: intuitionnistic logic (classical, linear, higher-order, constraints ...)

Deduction System : sequents calculus (LJ)

- A deduction rule:

$$\frac{\Gamma \vdash A \qquad \Gamma \vdash B}{\Gamma \vdash A \wedge B}$$

- right and left rules

$$\frac{}{\Gamma, A \vdash A}\text{axiom} \qquad \frac{\Gamma, A \vdash B \quad \Gamma \vdash A}{\Gamma \vdash B}\text{cut}$$

$$\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} \wedge\text{-r} \qquad \frac{\Gamma, A, B \vdash C}{\Gamma, A \wedge B \vdash C} \wedge\text{-l}$$

$$\frac{\Gamma, \forall x A[x], A[t] \vdash B}{\Gamma, \forall x A[x] \vdash B}\forall\text{-g, any } t \qquad \frac{\Gamma \vdash A[x]}{\Gamma \vdash \forall x A[x]}\forall\text{-r, } x \text{ free}$$

Example: 1

$$\forall x P(x) \vdash P(0) \wedge P(1)$$

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$$\frac{\forall x P(x) \vdash P(0) \qquad \forall x P(x) \vdash P(1)}{\forall x P(x) \vdash P(0) \land P(1)} \land\text{-r}$$

Example: 1

$$\dfrac{\forall\text{-l}\ \dfrac{\forall x P(x), P(0) \vdash P(0)}{\forall x P(x) \vdash P(0)} \qquad \dfrac{\forall x P(x), P(1) \vdash P(1)}{\forall x P(x) \vdash P(1)}\ \forall\text{-l}}{\forall x P(x) \vdash P(0) \wedge P(1)}\ \wedge\text{-r}$$

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Example: 2

$$\forall x P(x) \vdash P(0) \wedge P(1)$$

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$$\dfrac{\dfrac{\forall x P(x), P(1), P(0) \vdash P(0) \wedge P(1)}{\forall x P(x), P(0) \vdash P(0) \wedge P(1)}\ \forall\text{-l}}{\forall x P(x) \vdash P(0) \wedge P(1)}\ \forall\text{-l}$$

Example: 2

$$
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- the first rule is not always “don’t care”: free variable condition.

Axioms vs. rewriting

Axioms	Rewriting
$x + S(y) = S(x + y)$	$x + S(y) \rightarrow S(x + y)$
$x + 0 = x$	$x + 0 \rightarrow x$
$x * 0 = 0$	$x * 0 \rightarrow 0$
$x * S(y) = x + x * y$	$x * S(y) \rightarrow x + x * y$
$(x * y = 0) \Leftrightarrow (x = 0 \vee y = 0)$	$(x * y = 0) \rightarrow (x = 0 \vee y = 0)$
$\dfrac{\dfrac{\vdots}{\mathcal{T} \vdash 2 * 2 = 4}}{\mathcal{T} \vdash \exists x(2 * x = 4)}$	$\dfrac{\overline{\vdash 4 = 4}}{\vdash \exists x(2 * x = 4)}$

Deduction modulo: allowed rewriting

- General form (free variables are possible):

$$l \to r$$

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- we obtain a congruence modulo $\mathcal{R}$ (chosen set of rules): $\equiv$
- deduction rules transform as such:

$$\text{axiom}\ \frac{}{\Gamma, A \vdash A} \qquad \text{becomes} \qquad \frac{}{\Gamma, A \vdash B}\ \text{axiom},\ A \equiv B$$

Deduction modulo : sequent calculus modulo

$$\frac{}{\Gamma, A \vdash B}\text{axiom } A \equiv B \qquad \frac{\Gamma, A \vdash C \quad \Gamma \vdash B}{\Gamma \vdash C}\text{cut } A \equiv B$$

$$\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash C} \wedge\text{-r } A \wedge B \equiv C \qquad \frac{\Gamma, A, B \vdash C}{\Gamma, D \vdash C} \wedge\text{-l } A \wedge B \equiv D$$

$$\frac{\Gamma, B, A[t] \vdash C}{\Gamma, B \vdash C}\forall\text{-l } \forall x A[x] \equiv B \qquad \frac{\Gamma \vdash A[x]}{\Gamma \vdash B}\forall\text{-r}^* \ \forall x A[x] \equiv B$$

Example: 3

- consider the rewriting system $\mathcal{R}$:

$$\begin{array}{rcl} P(0) & \rightarrow & A \\ P(1) & \rightarrow & B \end{array}$$

$$\forall x P(x) \vdash A \wedge B$$

Example: 3

- consider the rewriting system $\mathcal{R}$:

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Cut rule: a detour

$$\frac{\Gamma, A \vdash B \quad \Gamma \vdash C}{\Gamma \vdash B} \text{ cut, } A \equiv C$$

- show $\Gamma \vdash A$
- show $\Gamma, A \vdash B$
- then, you have showed $\Gamma \vdash B$
- it is the application of a lemma.

Example: 4

- consider the rewriting system $\mathcal{R}$:

$$\begin{array}{rcl} P(0) & \rightarrow & A \\ P(1) & \rightarrow & B \end{array}$$

$$\forall x P(x) \vdash A \wedge B$$

Example: 4

- consider the rewriting system $\mathcal{R}$:

$$\begin{aligned} P(0) &\rightarrow A \\ P(1) &\rightarrow B \end{aligned}$$

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Example: 4

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Example: 4

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$$

- an unnecessary detour
- we could have cutted on any formula!

The cut rule: a detour

$$\frac{\Gamma, A \vdash B \quad \Gamma \vdash C}{\Gamma \vdash B} \text{ cut } A \equiv C$$

- we show $\Gamma, A \vdash B$ and $\Gamma \vdash A$
- then we have showed $\Gamma \vdash B$.
- lemma: the good way for a human being.
- in practice: not adapted for automatic demonstration.

  Nb: resolution method *do not* proceed by cuts !

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- eliminating cuts: a key result.

$$\Gamma \vdash A \quad \triangleright \quad \Gamma \vdash_{cf} A$$

- two main paths towards:
  - proof normalization (syntactic).
  - semantical methods.

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- two main paths towards:
  - proof normalization (syntactic).
  - semantical methods.
- in deduction modulo: indecidable, need for general criterions on $\mathcal{R}$

The normalization method(s)

- Curry-Howard: proofs = programs
- formulas = types
- proof tree = typing tree
- at the heart of proof assistants (PVS, Coq, Isabelle, ...)
- when a program calculates, it performs a cut elimination procedure.

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- at the heart of proof assistants (PVS, Coq, Isabelle, ...)
- when a program calculates, it performs a cut elimination procedure.
- show that all typables function terminates.

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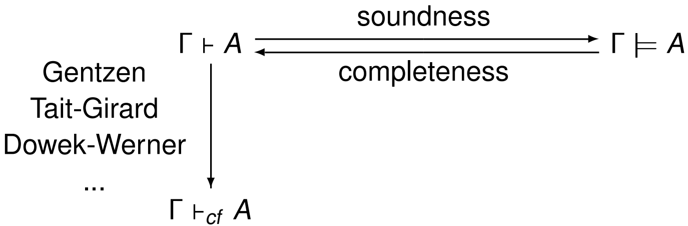
- we must have soundness/completeness wrt the semantic.
- there is links between both methods (last part of the talk).

The semantical method



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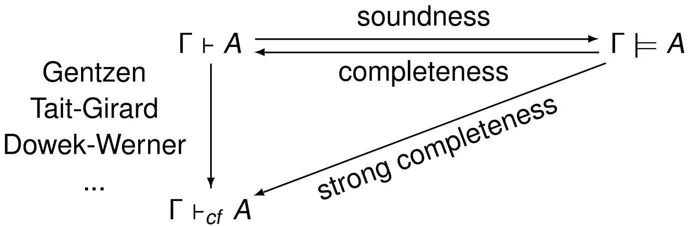

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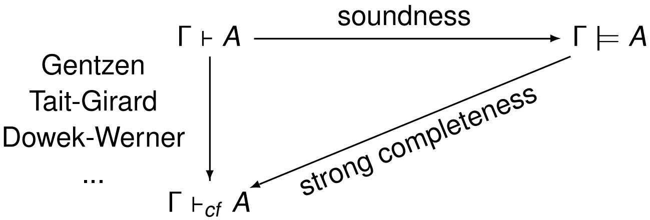

A semantic for deduction modulo

Two main semantics for intuitionistic logic:

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- Heyting algebras [Lipton,Okada]

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- K the set of worlds, partially ordered with $\leq$ (a "temporal relation": past, present, possible futures: partial information)

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- $D : \alpha \to \mathcal{S}et$ a monotone function (interpretation domain for terms).

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- K the set of worlds, partially ordered with $\leq$ (a "temporal relation": past, present, possible futures: partial information)
- $D : \alpha \rightarrow \mathcal{S}et$ a monotone function (interpretation domain for terms).
- $\Vdash$ is a relation between worlds and formulas, verifiying:

A semantic for deduction modulo

- P atomic: if $\alpha \leq \beta$ and $\alpha \Vdash P$, then $\beta \Vdash P$.
- $\alpha \Vdash A \Rightarrow B$ iff for any $\beta \geq \alpha$, when $\beta \Vdash A$ then $\beta \Vdash B$.
- $\alpha \Vdash A \vee B$ iff $\alpha \Vdash A$ or $\alpha \Vdash B$.

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- $\alpha \Vdash A \vee B$ iff $\alpha \Vdash A$ or $\alpha \Vdash B$.
- Additional constraint in deduction modulo:

$$A \equiv B \quad \text{implies} \quad \alpha \Vdash A \Leftrightarrow \alpha \Vdash B$$

Kripke structures at work

- $A \vee (\neg A)$ is well-known not to be valid in intuitionistic logic.
- we build a structure that is invalidating this formula. Note: at least two worlds (single world = boolean model).
- $\neg A = A \Rightarrow \bot$

$$\beta \Vdash A$$

$$|$$

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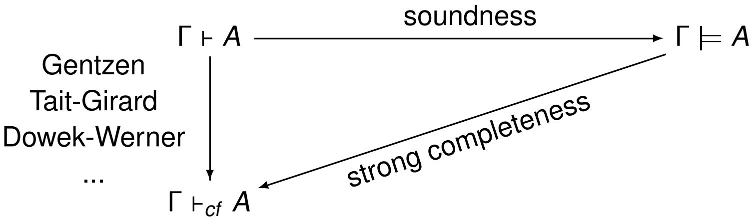
$$\alpha \Vdash \emptyset \text{ and } \alpha \nVdash A, \neg A, A \vee \neg A$$

Constructive proof: the algorithm behind



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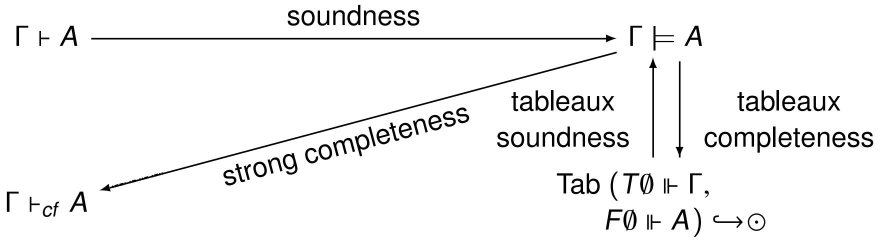

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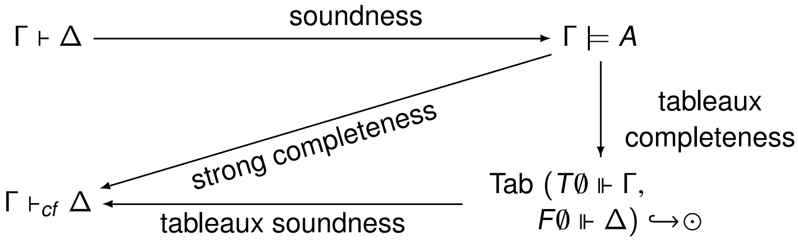

The tableau method

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- some rules:

$$\begin{array}{c} Tp \Vdash A \vee B \\ \diagup \quad \diagdown \\ Tp \Vdash A \qquad Tp \Vdash B \end{array} \qquad\qquad \begin{array}{c} Fp \Vdash A \vee B \\ | \\ Fp \Vdash A \\ | \\ Fp \Vdash B \end{array}$$

$$\begin{array}{c} Tp \Vdash A \Rightarrow B \\ \diagup \quad \diagdown \\ Tq \Vdash B \qquad Fq \Vdash A \end{array} \qquad\qquad \begin{array}{c} Fp \Vdash A \Rightarrow B \\ | \\ Tq \Vdash A \\ | \\ Fq \Vdash B \end{array}$$

with proviso on q

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with proviso on q

- in deduction modulo: allow rewrite rules, define a new systematic research algorithm with $\mathcal{R}$.

Tableau: example 1

- We want to show "$A \vee B \vdash C \Rightarrow A$"
- tranlsation in tableau language: there is NO (node of no) Kripke structure satisfying $A \vee B$ without satisfying also $C \Rightarrow A$. Let's see if the counter-model search fails or not.
- We choose as usual sequences of integers for the set of worlds (partial order: prefix).

$T\emptyset \Vdash A \vee B, F\emptyset \Vdash C \Rightarrow A$

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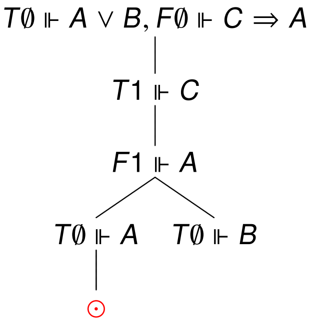

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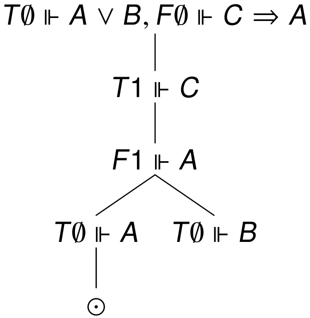

Tableau: example 2

- We want to show "$\vdash (A \Rightarrow B) \Rightarrow (A \Rightarrow B)$"

$F_{\varnothing} \Vdash (A \Rightarrow B) \Rightarrow A \Rightarrow B$

Tableau: example 2

$$F_{\emptyset} \Vdash (A \Rightarrow B) \Rightarrow A \Rightarrow B$$

$$T_1 \Vdash (A \Rightarrow B)$$

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Tableau: example 2

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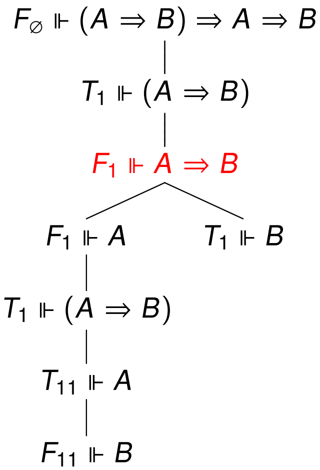
$$F_1 \Vdash A \qquad\qquad T_1 \Vdash B$$

$$T_1 \Vdash (A \Rightarrow B)$$

Tableau: example 2



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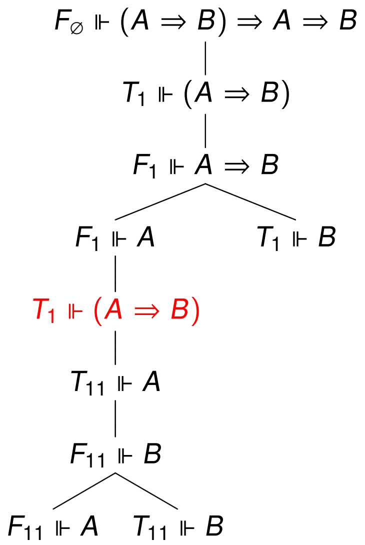
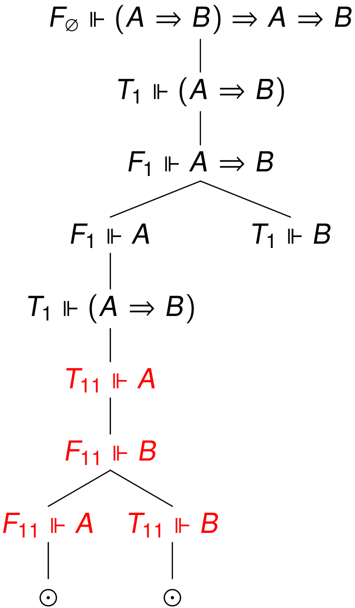


Tableau: example 2



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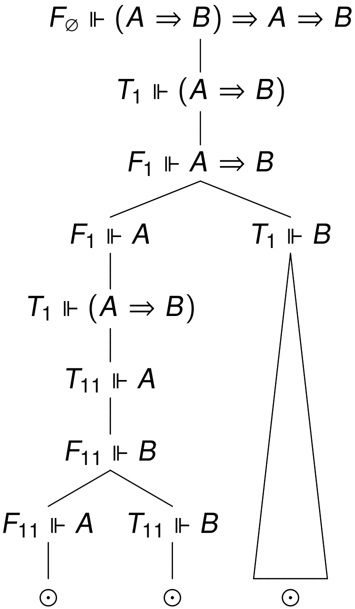

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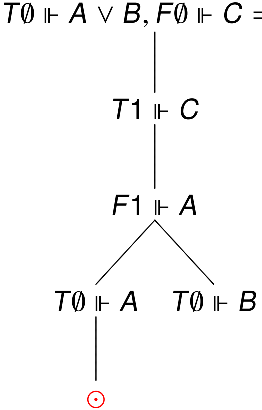
  - deduction modulo: it has also to be a model of the rewrite rules $\mathcal{R}$.
- constructive point of view: if there is no counter-model, does the method terminate? (KS definition is modified)

Remember the tableau for $A \vee B \vdash C \Rightarrow A$:



- the right path generates counter model.
- the nerve: the atomic formulas each world entails (forces), extension by induction.

Conditions on rewrite rules

Providing the confluence of the rewrite system $\mathcal{R}$, and for:

- an order condition: $\succ$, well-founded, having the subformula property, and such that $P \rightarrow^* Q$ implies $P \succ Q$.

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$\Gamma \vdash \Delta \xrightarrow{\text{soundness}} \Gamma \models \Delta$

$\Gamma \models \Delta \xrightarrow{\text{tableaux completeness}} \text{Tab}\,(T\emptyset \Vdash \Gamma, F\emptyset \Vdash \Delta) \hookrightarrow \odot$

$\Gamma \vdash_{cf} \Delta$

Tableaux soundness

We show the following theorem:

Theorem

If a tableau starting with $T\emptyset \Vdash \Gamma, F\emptyset \Vdash P$ is closed, then we can transform it into a proof of $\Gamma \vdash_{cf} P$.

- intuitionistic diffculty: in a tableau, there might be more than one "non true" formula:

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  (similar to "multi succedent intuitionistic sequent calculus").
- easy with cut, hard without.

Normalization (in a nutshell)

Curry-Howard correspondence

- Notation for proofs:

$$\frac{\Gamma, x : A \vdash \pi : B}{\Gamma \vdash \lambda x.\pi : A \Rightarrow B} \qquad \frac{\Gamma \vdash \pi' : A \qquad \Gamma \vdash \pi : A \Rightarrow B}{\Gamma \vdash (\pi\pi') : B}$$

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- cut elimination is a **process**, similar to function execution.
- aim: show that every proof normalizes: then the cut elimination process terminates.

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- our aim: to each formula A, find a candidate $[\![A]\!]$. Show that if $\Gamma \vdash \pi : A$ then $\pi \in [\![A]\!]$.
- in deduction modulo, if $A \equiv B$, additional constraint:

$$[\![A]\!] = [\![B]\!]$$

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- such methods are defined in deduction modulo (Heyting arithmetic, higher-order logic, Zermelo’s set theory, ...)

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- the sets of candidates have a structure: pseudo Heyting algebras [Dowek].

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$$a \leq a \cup b \quad b \leq a \cup b \quad a \leq c \text{ and } b \leq c \text{ implies } a \cup b \leq c$$

- think about $\mathbb{R}$ and closed sets (infinite l.u.b. is not infinite union)

pseudo-Heyting algebras

- a universe Ω
- a pseudo order: $a \leq b$ and $b \leq a$ with $a \neq b$ possible.
- operations on it: lowest upper bound (join: $\cup$ – pseudo union), greatest lower bound (meet: $\cap$ – intersection).

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- but ... Heyting algebras used for semantical cut elimination.

The link: fibring

define

$$[A] = [\![A]\!] \lhd A = \{\Gamma \mid \Gamma \vdash \pi : A, \pi \in [\![A]\!]\}$$

- **weak** definition: for *some* π only.
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- Takahashi, Prawitz, Schütte, higher-order V-complexes (extended).

Computational content: what kind of algorithm ?

Let's consider the rule:

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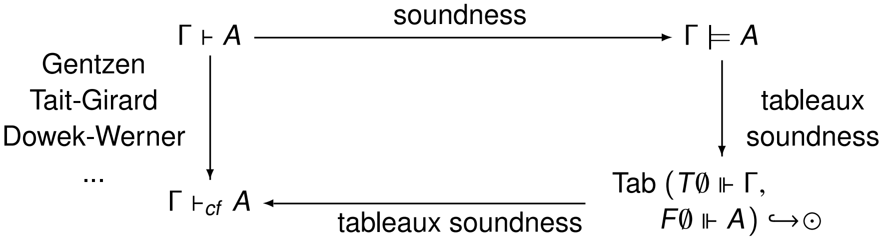
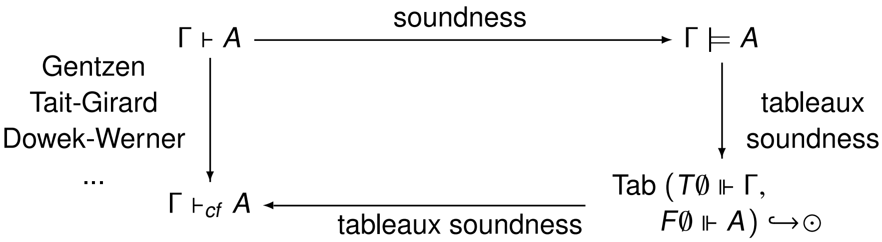
- this can not be a normalization algorithm.
- it is more or less the tableau method described in the first part.



- This diagram does not commute.



- This diagram does not commute.
- But: normalization methods “generate” a certain kind of semantical cut elimination proof: normalization by evaluation (**weak** fibring).

Further work

- there is normalization by evaluation work, but in a Kripke style: links with both works ?
- do the candidates always have a “pseudo-” structure ?
- realizing rewrite rule not with $\lambda x.x$ (not silently), could recover normalization.